

Mesh Partitioning for Parallel CFD Application on a Grid

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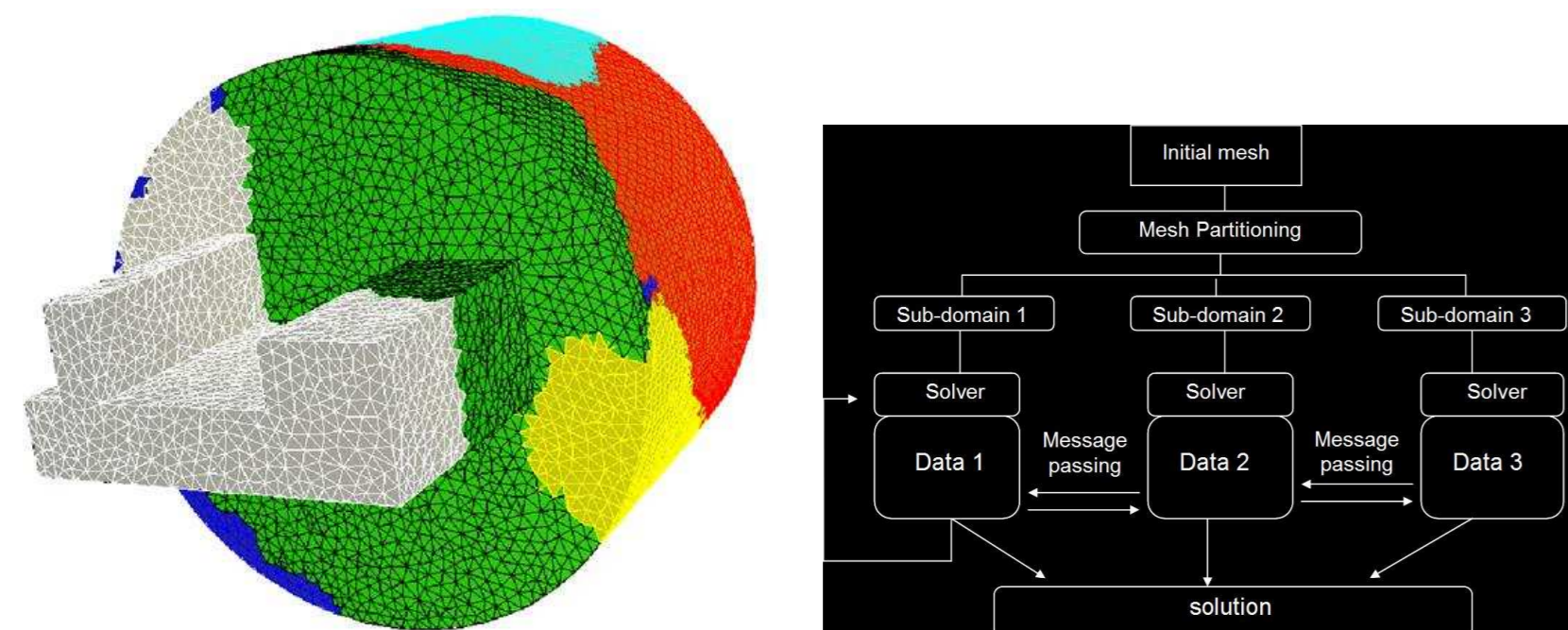
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1 CFD and parallelism SPMD model



2 Overview

- Mesh partitioning = efficiency of parallel CFD application
- Homogeneous partitioning : load balancing
- Existing schemes:
 - ParMetis
 - Jostle
 - Chaco
- But, grid architectures:
 - heterogeneity in CPU
 - heterogeneity in network links

3 Mesh/Architecture models

- Mesh application graph:

$W(V, E)$: Workload graph

$w(v)$: The weight represents the work load associated to this node

$w(\{u, v\})$: The weight represents data dependency between two nodes

- **Architecture graph:** $A(P, E)$: Complete graph

S_p : Processing power of p

v_{pq} : Bandwidth between p and q

4 Example of architecture: The MecaGrid graph

Mecagrid: project started 11/2002

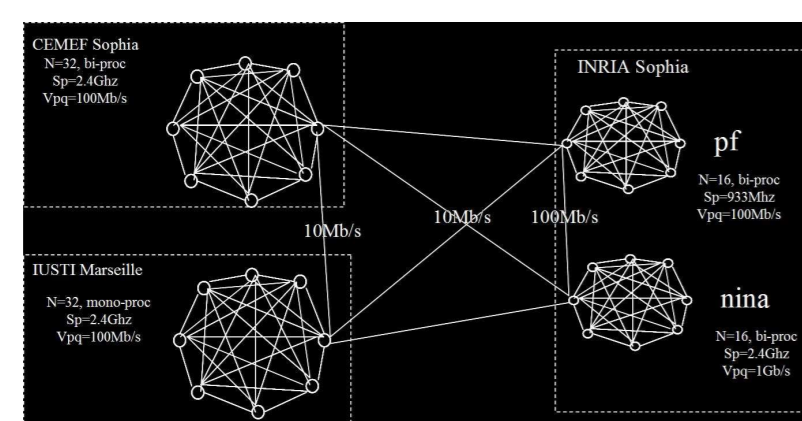
Connects 3 sites in the PACA (Provence Alpes Côte d'Azur) region

Performs experiments in grid computing applied to multimaterial fluid dynamics

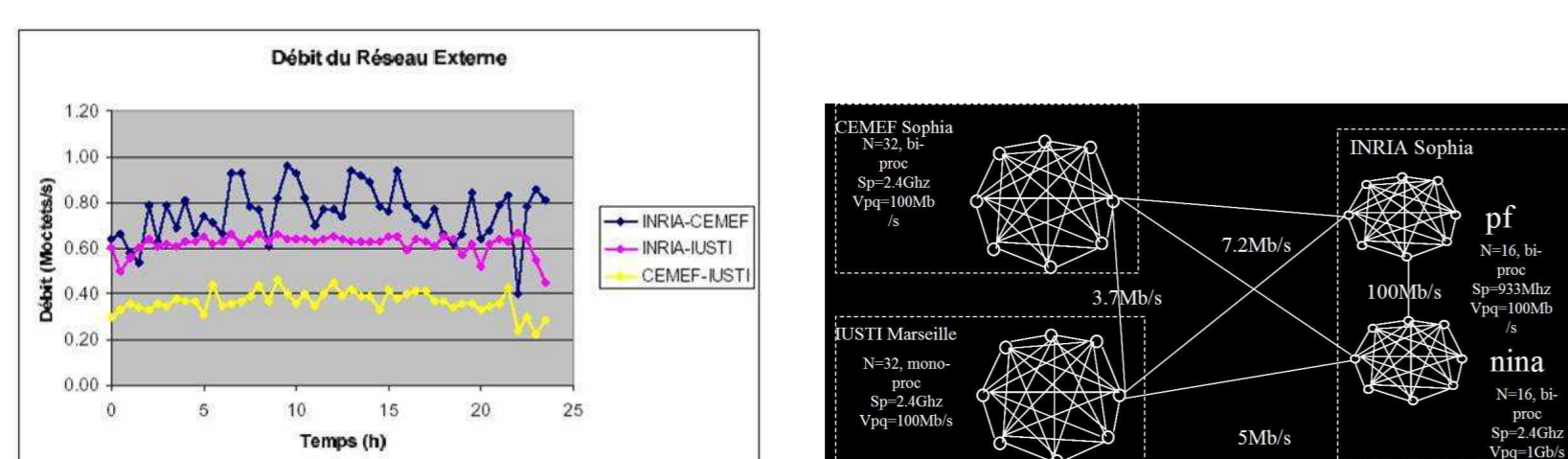


5 The MecaGrid: heterogeneous architecture of 162 processors

a The MecaGrid: Graph architecture



b The MecaGrid: measured performances



6 Heterogeneous mesh partitioning

The mapping problem: Find a mesh partition that minimize CPU time.

Homogeneous (cluster architecture): Standard load balancing

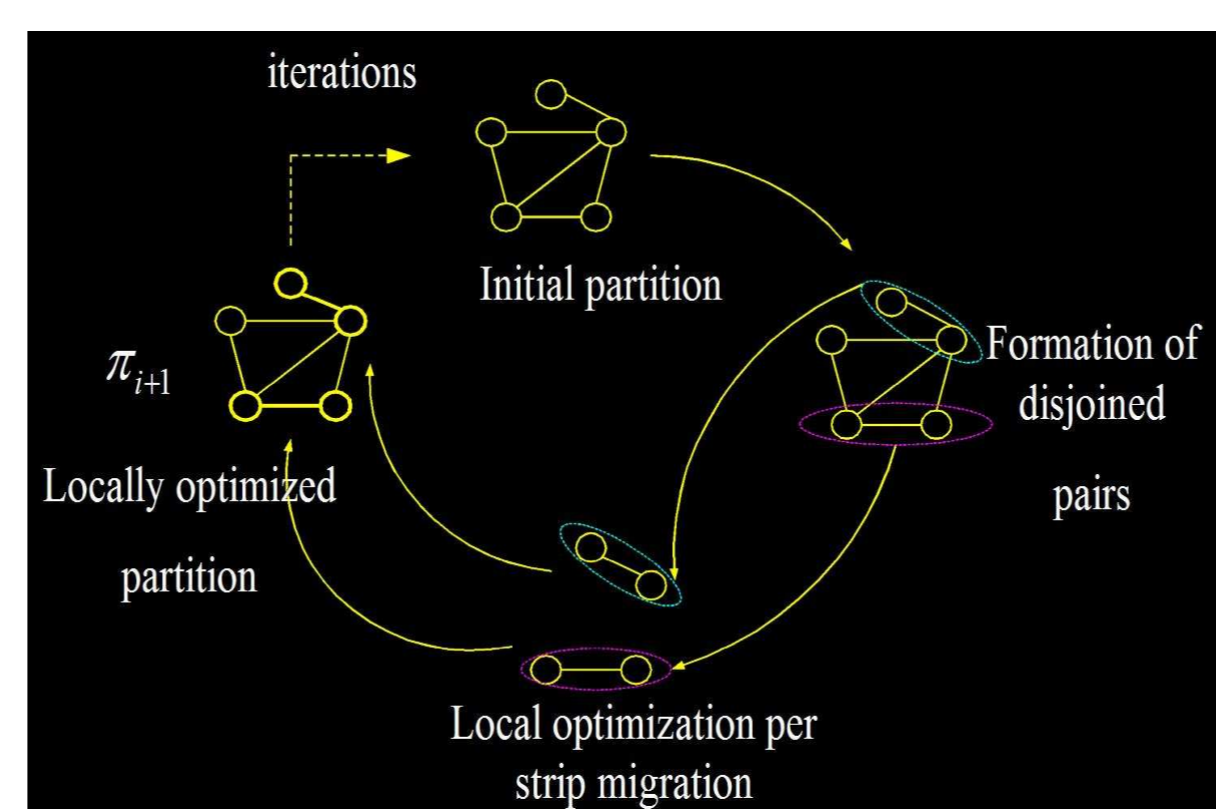
Heterogeneous (Grid):

$$F(W, A, m) = \max_{p \in V(A)} (t_p) + \max_{p \in V(A)} (C_p)$$

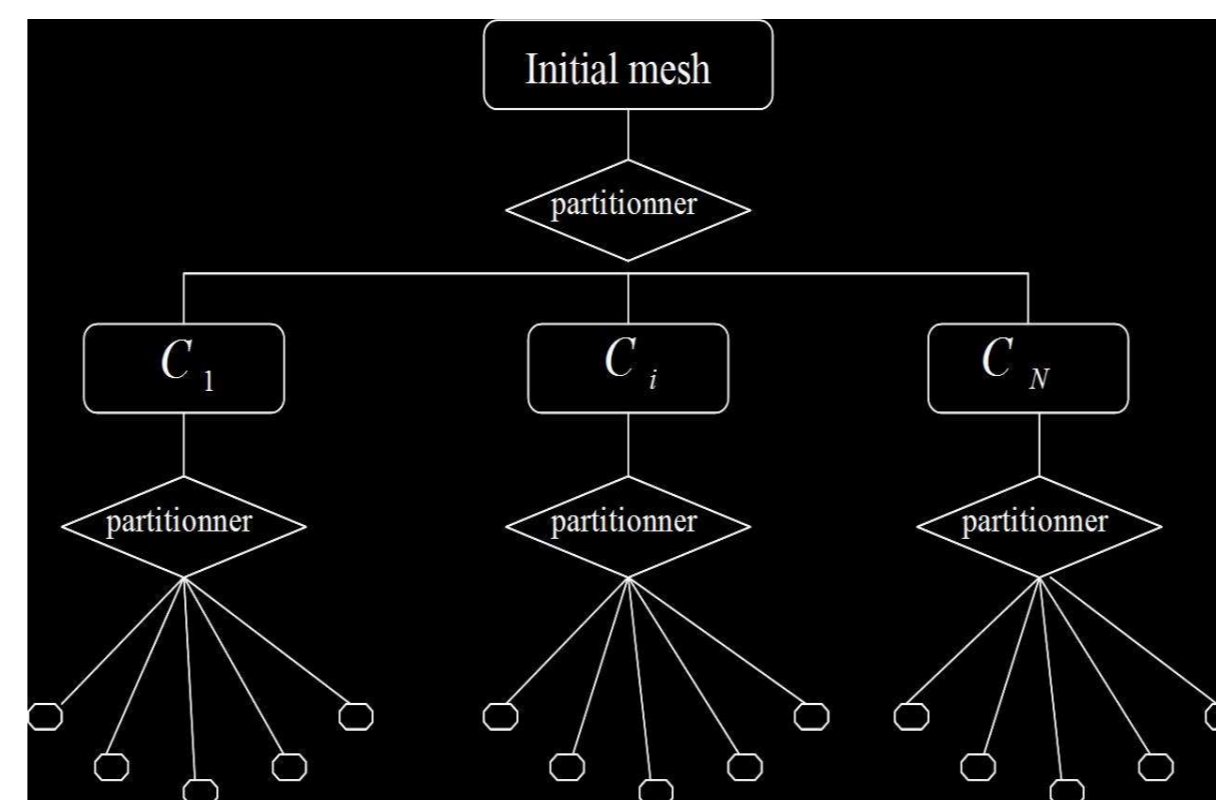
$$t_p = \frac{C(p, m)}{S_p} \quad \text{where} \quad C(p, m) = \sum_{\substack{v \in V(S) \\ m(v)=p}} w(v)$$

$$c_p = \sum_{\substack{i \in V(S) \\ p \neq i}} c_{pi} \quad \text{where} \quad c_{pi} = \frac{C(\{p, i\}, m)}{v_{pi}}$$

7 Parallel partitioning algorithm



8 Hierarchical mesh partitioner



9 Application examples and validation

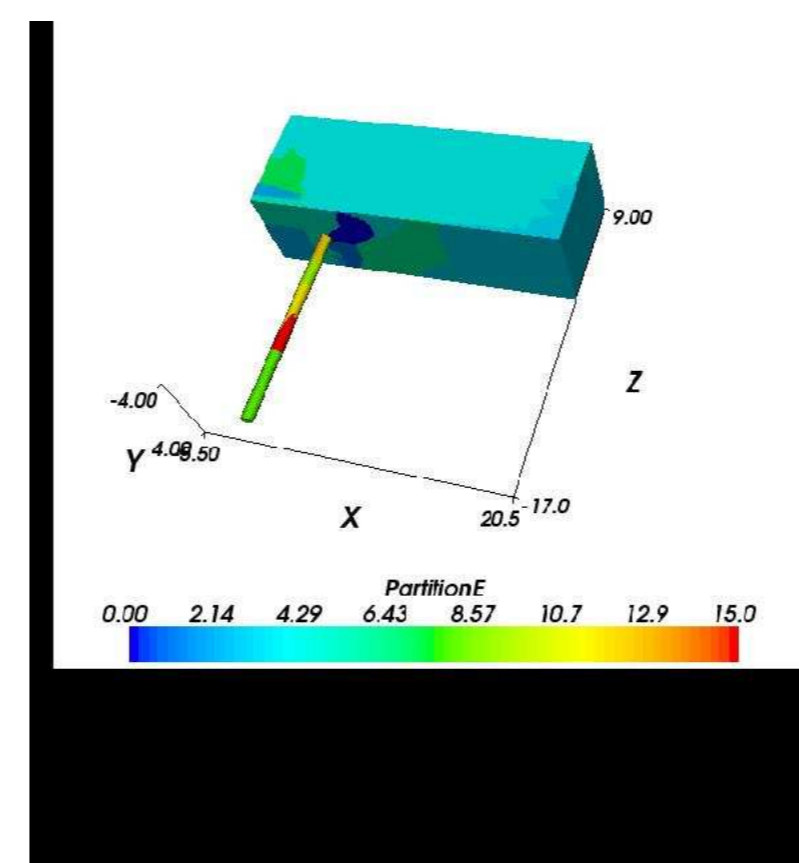
Tools: CIMlib library of CEMEF: a C++/MPI finite element library solving multimaterial flows (H. Dignonnet, O. Basset, T. Coupez, H. Guillard)

1 Exemple 1:

a Problem to solve:

$$(1) \quad \begin{cases} \nabla \cdot (2\eta \epsilon(v)) - \nabla p = 0 \\ \nabla \cdot v = 0 \end{cases}$$

b Mesh: 3D mesh with 398K nodes and 3.8M elements



c Results on the INRIA clusters

Sop1-Sop2	CPU time	Sop1-Sop2	CPU time
Resolution (s)	339.97	Resolution (s)	174.22
Assembling (s)	9.40651	Assembling (s)	6.61217
Total (s)	349.42	Total (s)	180.86

Homogeneous partition

Optimized partition

Gain: 50%

d results on the INRIA-MARSEILLE clusters

16MRS-16Sop2	CPU time	16MRS-16Sop2	CPU time
Resolution (s)	561.278	Resolution (s)	137.924
Assembling (s)	18.227	Assembling (s)	5.735
Total (s)	579.505	Total (s)	143.66

Homogeneous partition

Optimized partition

Gain: 75%

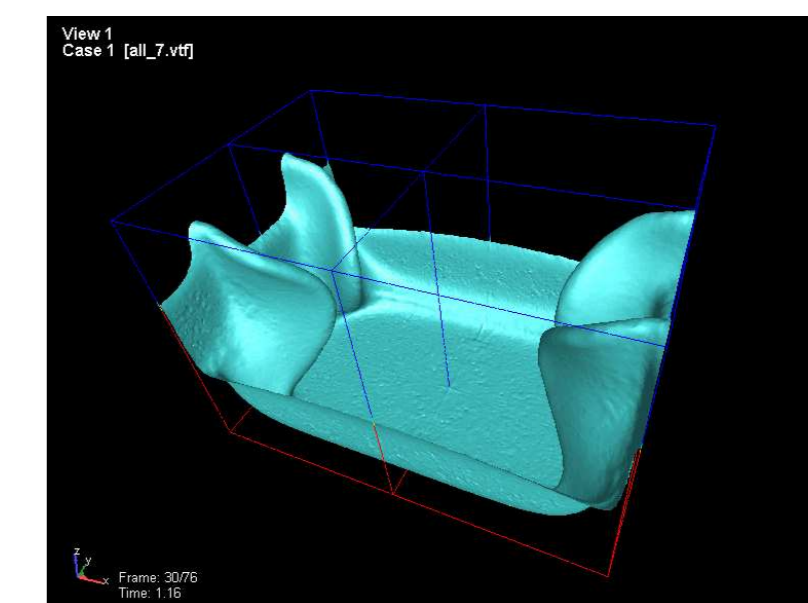
2 Example 2:

a Multi-fluid unsteady incompressible NS equations coupled with transport equation

$$\begin{cases} \rho(\alpha) \frac{dv}{dt} - \nabla \cdot (2\eta \epsilon(v)) + \nabla p = \rho(\alpha) g \\ \nabla \cdot v = 0 \\ \frac{\partial \alpha}{\partial t} + \nabla \cdot \alpha v = 0 \end{cases}$$

Viscosity & Density depend on a phase variable α

b Mesh: 3D mesh with 1.5M nodes and 8.7M elements



c NS and levelset + continuous Galerkin + Hamilton-Jacobi reinitialisation

Mesh: 1.5M nodes and 8.7M elements 600 time steps
2 linear systems per increment of 6 millions of unknowns and of 1.5 millions respectively: a total of 4 billion and 500 millions of unknowns

Clusters(39 procs used)	CPU time
Sop1-Sop2-MRS(hom)	5 days and 6 hours
Sop1-Sop2-MRS(het)	3 days and 1 hour

Gain: 45%

Conclusion

- Presentation of a new parallel partitioning scheme for unstructured meshes on heterogeneous architectures.
- Real life validation of the method with FE code executed on the real life grid.
- Gain of more than of 75% on CPU time computation

