

SPLASH

Securing Adhoc Networks

Eurecom- INRIA - UC Irvine

2003-2006

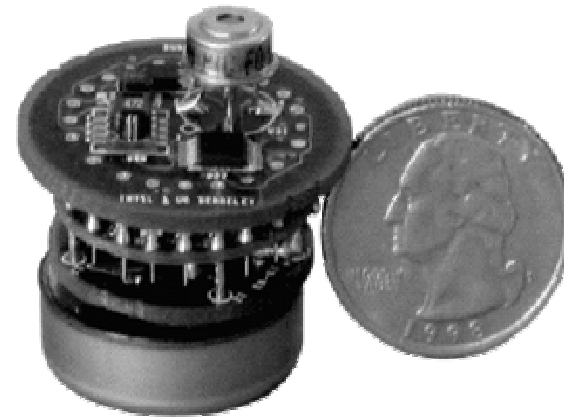
www.inrialpes.fr/planete/splash.html

Claude Castelluccia

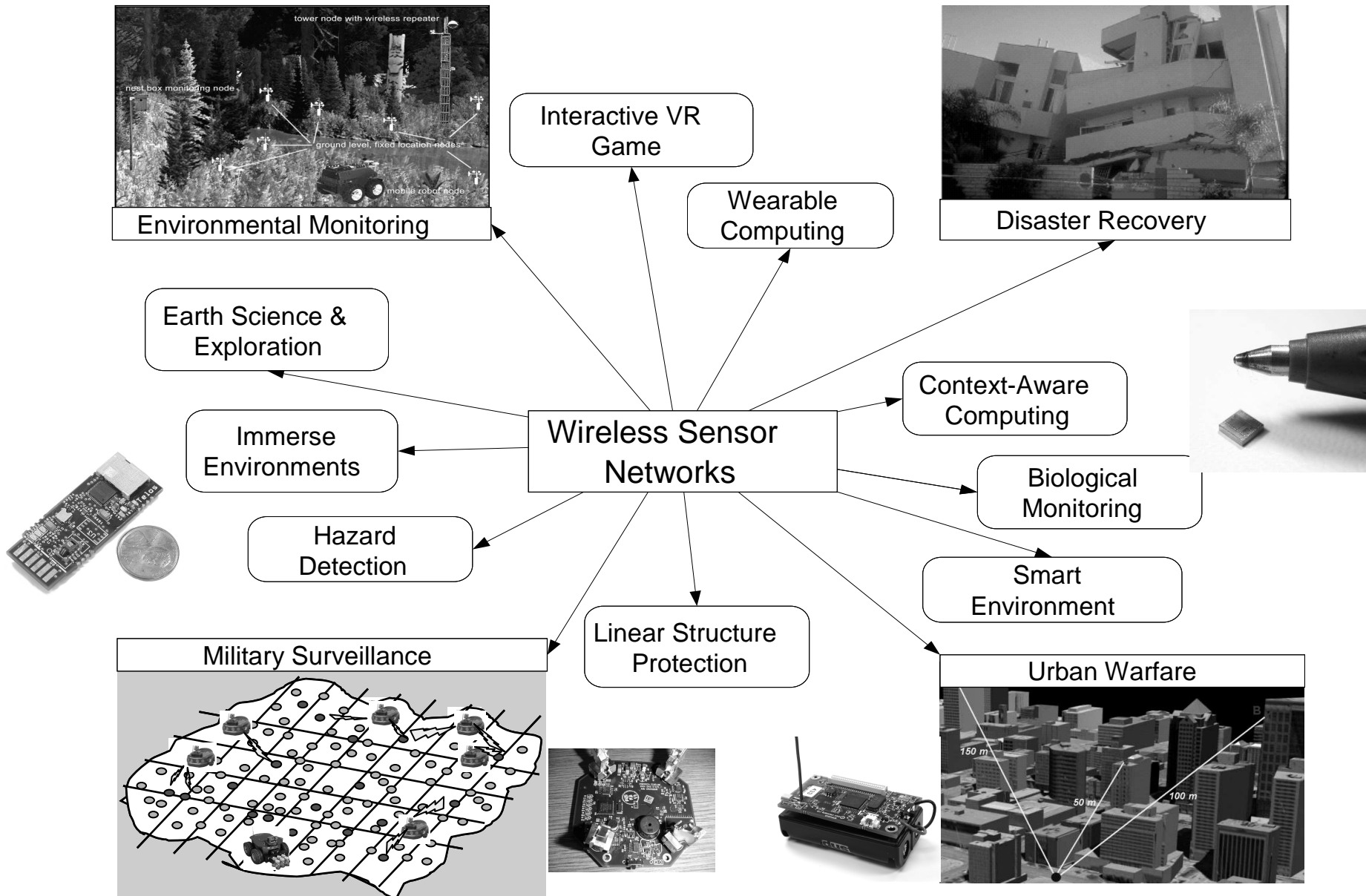
Nov. 2005

The SPLASH Project: *Securing Adhoc Networks*

- Adhoc Networks
 - MANET: Mobile Adhoc networks (distributed, no central authority)
 - Wireless Sensor Networks (constraint devices/ one central authority)
- Characteristics
 - Limited CPU/Battery: need light-weight solutions
 - Fully distributed: no Central authority can be assumed



Application Spectrum



MANET Security Requirements

Wireless & Mobile

- Limited Energy
- Lack of physical security

• Cooperation Enforcement

Ad Hoc

- Lack of (or limited) infrastructure
- Lack of a priori trust

• Secure Routing

• Key management

*[Recent security solutions for mobile ad hoc networks
In "Ad Hoc Networks" IEEE Press - Wiley Ed]*

The SPLASH Project:

Securing Adhoc Networks

- Some of our Research Challenges
 - Efficient and Infrastructure-less Key establishment/pairing
 - Distributed Access Control
 - Collaboration Enforcement
 - Securing Routing Protocols
 - Secure Aggregation
 - Privacy

One example of our results:

Shake Them Up!

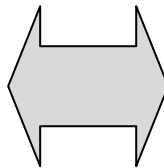
ACM Mobysis 2005, Seattle, USA

Shake Them Up!

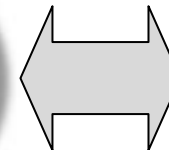
- In many wireless applications, you need to pair, i.e. establish a shared secret on-the-fly between devices.
- Some devices, such as sensors, have *very limited* CPU, memory and power!
- Standard methods such as the DH key exchange are excluded!
- Examples:



Wireless
Thermometer



PDA



Current Solutions

- Public Key Cryptography-based schemes
 - Rely on PK key exchange protocols such as RSA or DH
 - They require to perform CPU-intensive operations such as modular exponentiations with large numbers!
 - Too expensive for sensor devices!
- PIN-based schemes (for ex. Bluetooth)
 - Key derived from a PIN number
 - PIN number is typically entered via an out-of-band channel such as a keyboard.
 - Computationally efficient.
 - ...but requires a physical user interface (keyboard)
...and most sensors do not have a keyboard \perp !
 - Security is pretty weak since it depends on the PIN....

Current Solutions (2)

- Physical Contact (imprinting)
 - Stajano and Anderson proposed to establish a key via physical contact by linking devices with a wire....
 - Not always practical and requires additional hardware..
- InfraRed channel
 - IR is difficult to intercept since requires line-of-sight links.
 - But most sensors do not have IR interface!
- Faraday Cage
 - Devices could be placed into a Faraday cage
 - It is clearly impractical to ask users to lug around a metal box ;-)

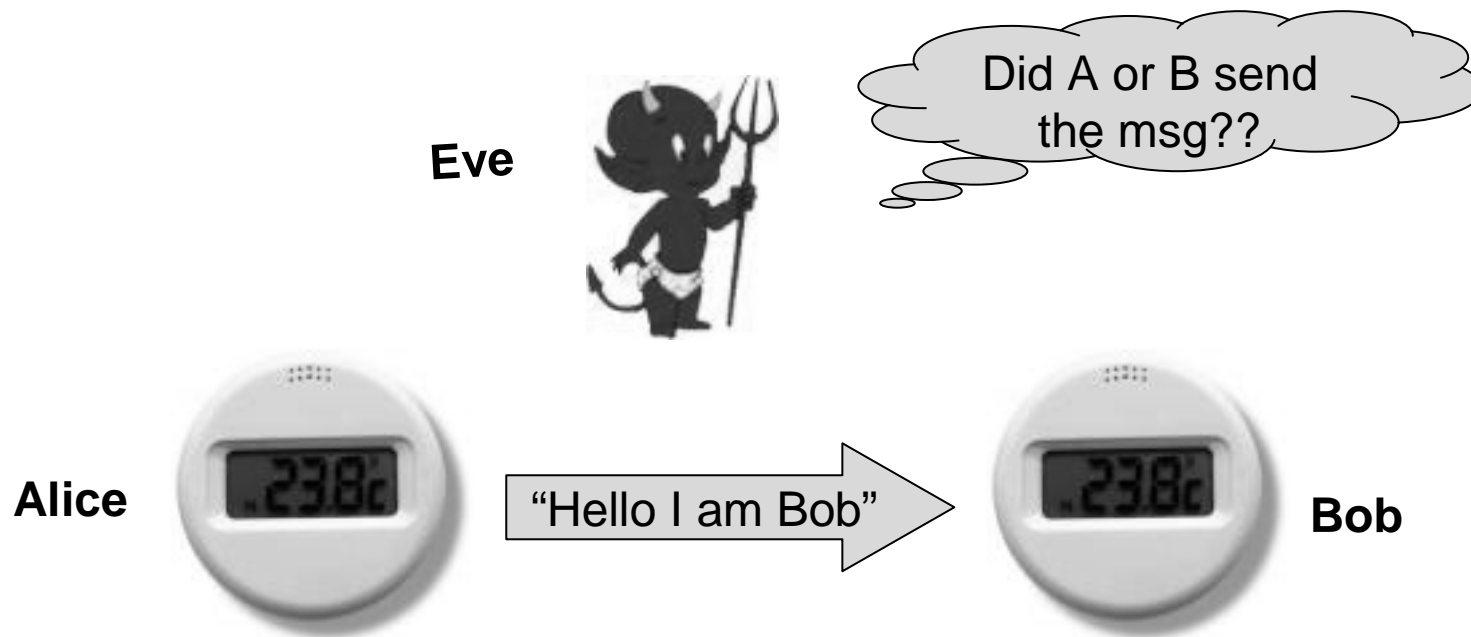
Our Goals

- Design a secure pairing protocols that:
 - Does not rely on PK cryptography
 - Does not rely on pre-configured information
 - Does not increase the complexity (and cost) of the sensors by requiring additional hardware such as a display, keyboard, IR channel...
 - Does not require special equipments (cable, faraday cage)
- Security Model
 - The protocol must ensure that active or passive attackers do not learn the exchanged key
 - It must provide some DoS protection, i.e. prevent an attacker from disrupting the key exchange and exhausting the devices' resources.

Our solution:

How to exchange one secret bit

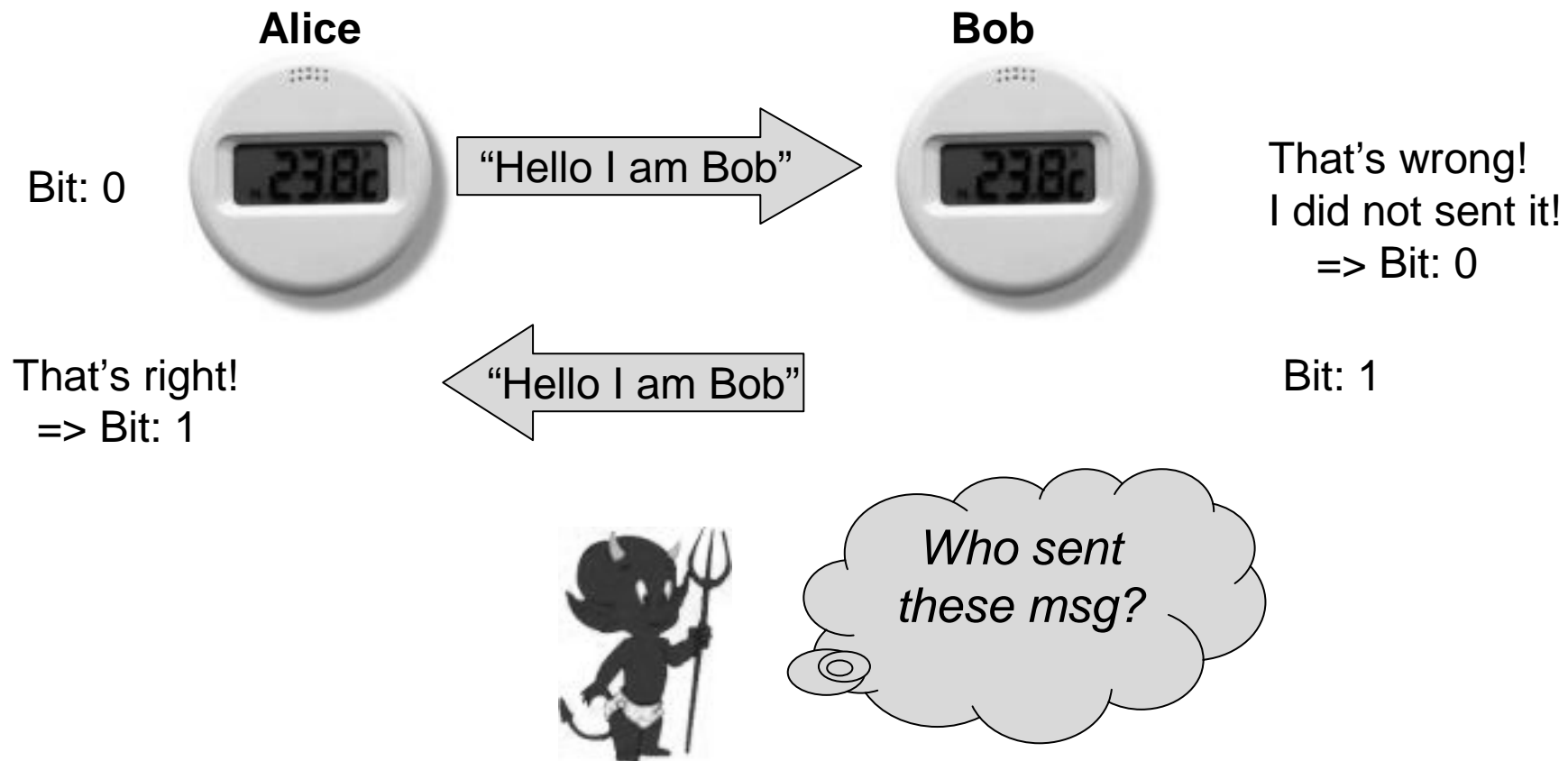
- Let's assume that Alice (A) and Bob (B) communicate over a wireless *anonymous* broadcast channel
 - Eve can read the exchanged packets
 - ...but can not identify the source of the packets.



Our solution:

How to exchange one secret bit (2)

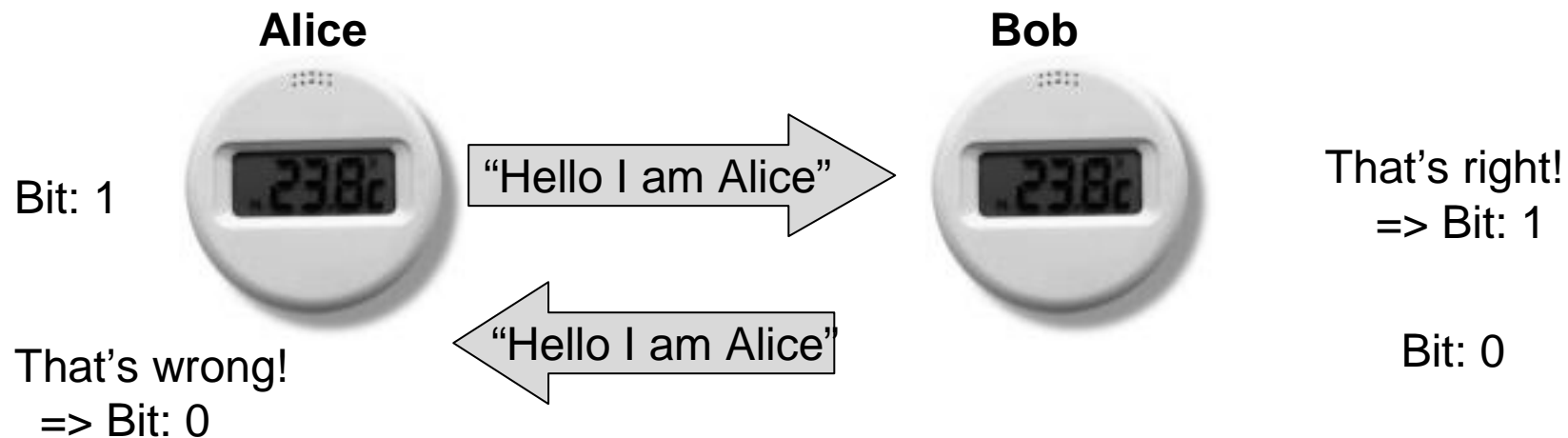
- Alice and Bob can then use the following algorithm:



Our solution:

How to exchange one secret bit (3)

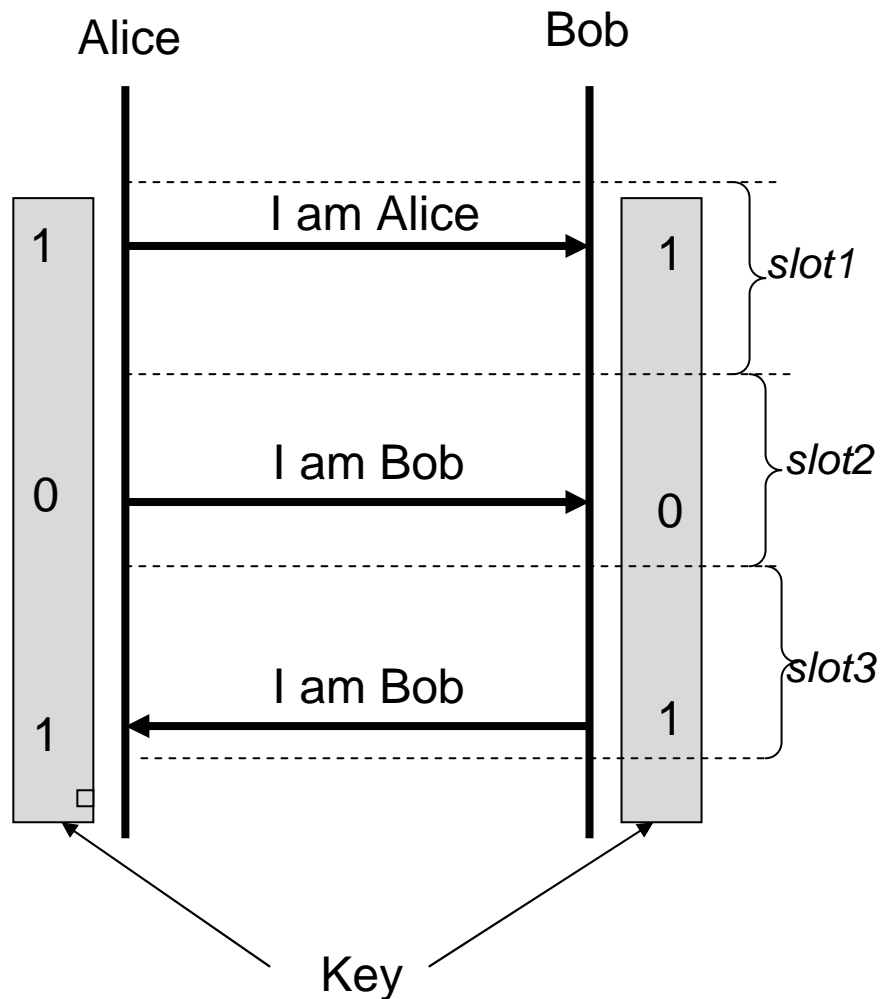
- Of course the protocol is symmetrical i.e. Alice can also send the bit "1" and Bob the bit "0"



Our solution:

How to exchange a N-bit secret

- We divide the time in N slots.
- During each slot, either A or B sends a message
- The transmission order is random so that Eve can not group the messages together and retrieve the key....!

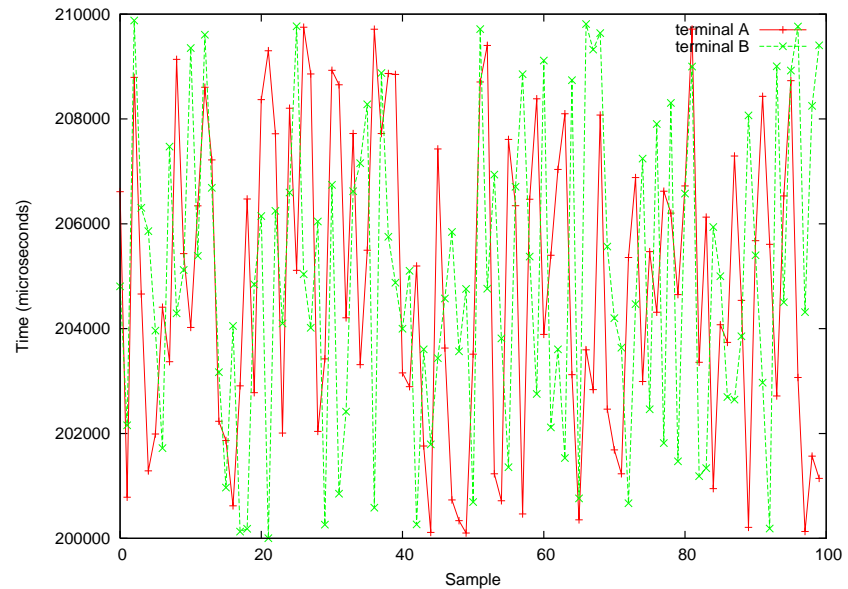


Wireless Anonymous Communication

- We assume source anonymity...
 - Can an 802.11-based system provide source anonymity?
- Eve can potentially identify the real source of the messages
 - Timing information
 - Reception Power
 - Frequency

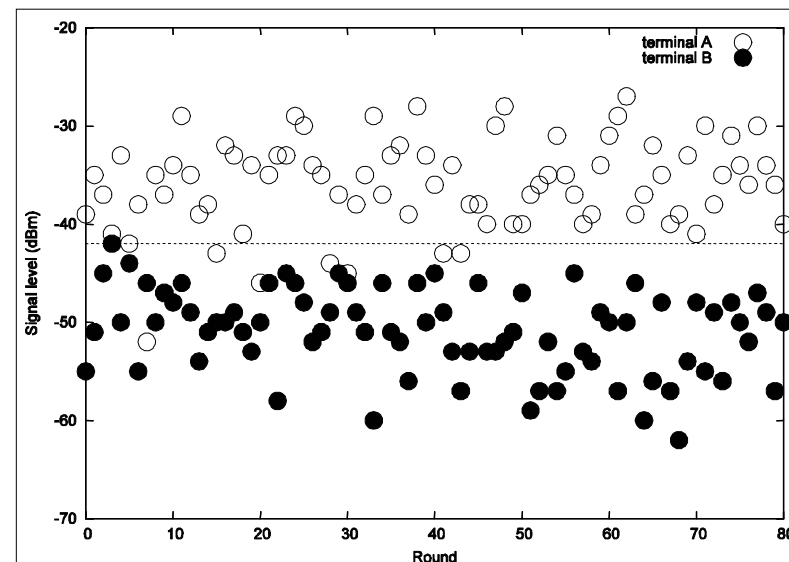
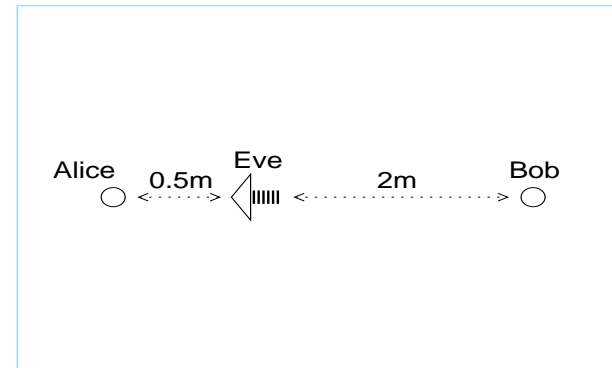
Wireless Anonymous Communication (2)

- Timing
 - This is quite trivial in TDMA based scheme since devices always transmit during their allocated slots
 - However Timing does not provide any information if a random access MAC protocol, such as CSMA, is used since each device access the channel at a random time!
- => Our protocol only works with CSMA-based technologies, such 802.11,802.15.4



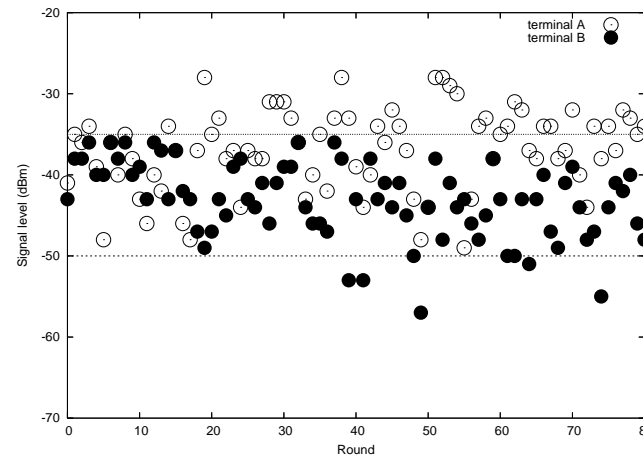
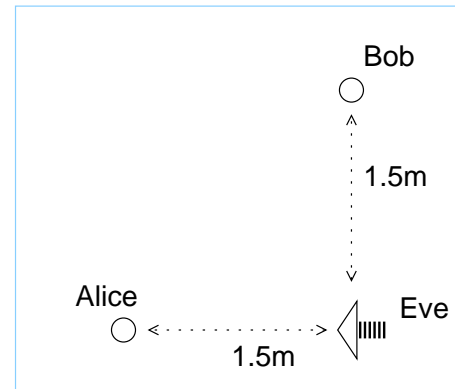
Wireless Anonymous Communication (3)

- Reception Power
 - The reception power is defined as
$$S_r = S_t \cdot G_t \cdot G_r \cdot K / d^2$$
 - If Eve is closer to Alice than Bob, she will receive Alice's message with a higher power!
 - Note: we assume A and B transmit at the same power level.



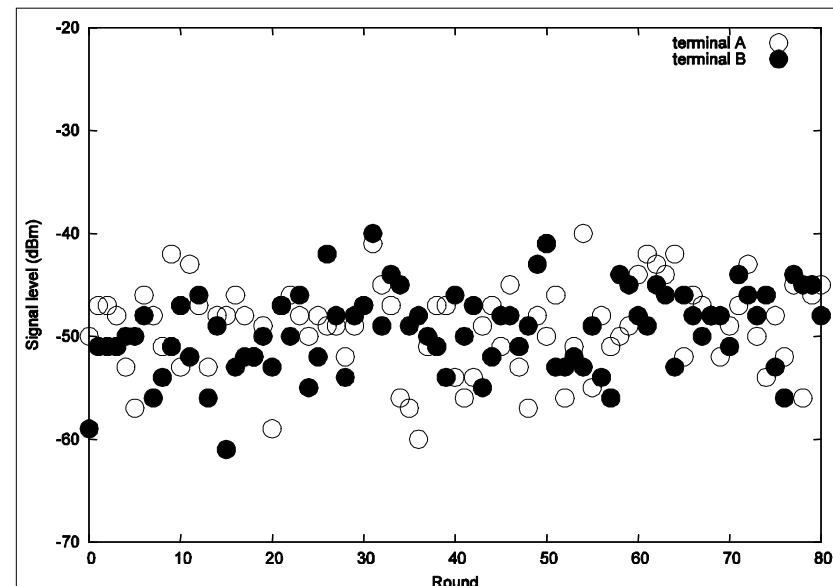
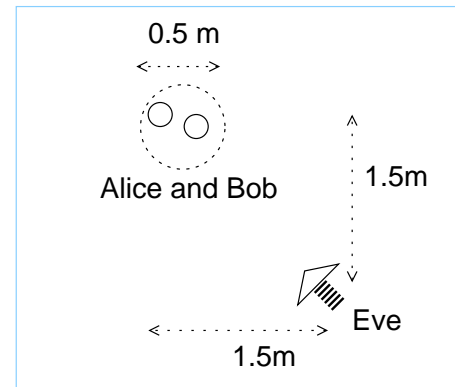
What can be done? (1)

- We can randomly change Alice and Bob's transmission power
 - Some bits will still be revealed
 - If Eve has a directional antenna she can aimed it at one of the devices!



What can be done? (2)

- We can bring the devices together and move them (shake them up) one around the other!
 - The reception power of A's and B's messages will be similar...
 - Eve cannot use a directional antenna since the devices are moving!
- In summary, shaking 2 devices prevents using power to identify source!



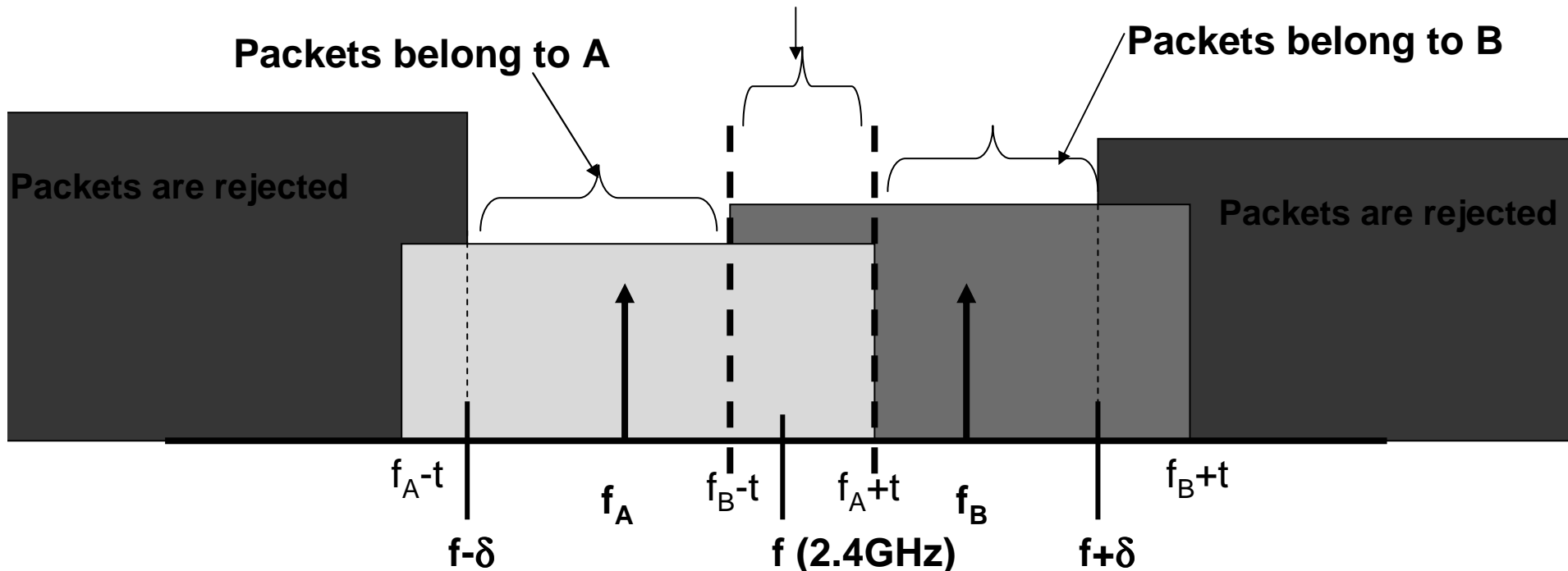
Frequency Fingerprinting

- Even though standards specify one frequency, each device uses a different frequency.
- This difference is due to the crystal oscillator and clock drift, resulting from aging, temperature and so on.
- Typically an error of 25ppm (parts per million) is allowed by the standard.
- So if the transmitting frequency is 2.4GHz, a frequency offset of up to 120kHz is allowed.
- Possibly, a (well-equipped) Eve can use this frequency difference to identify the source and retrieve the secret...

Frequency Fingerprinting (2)

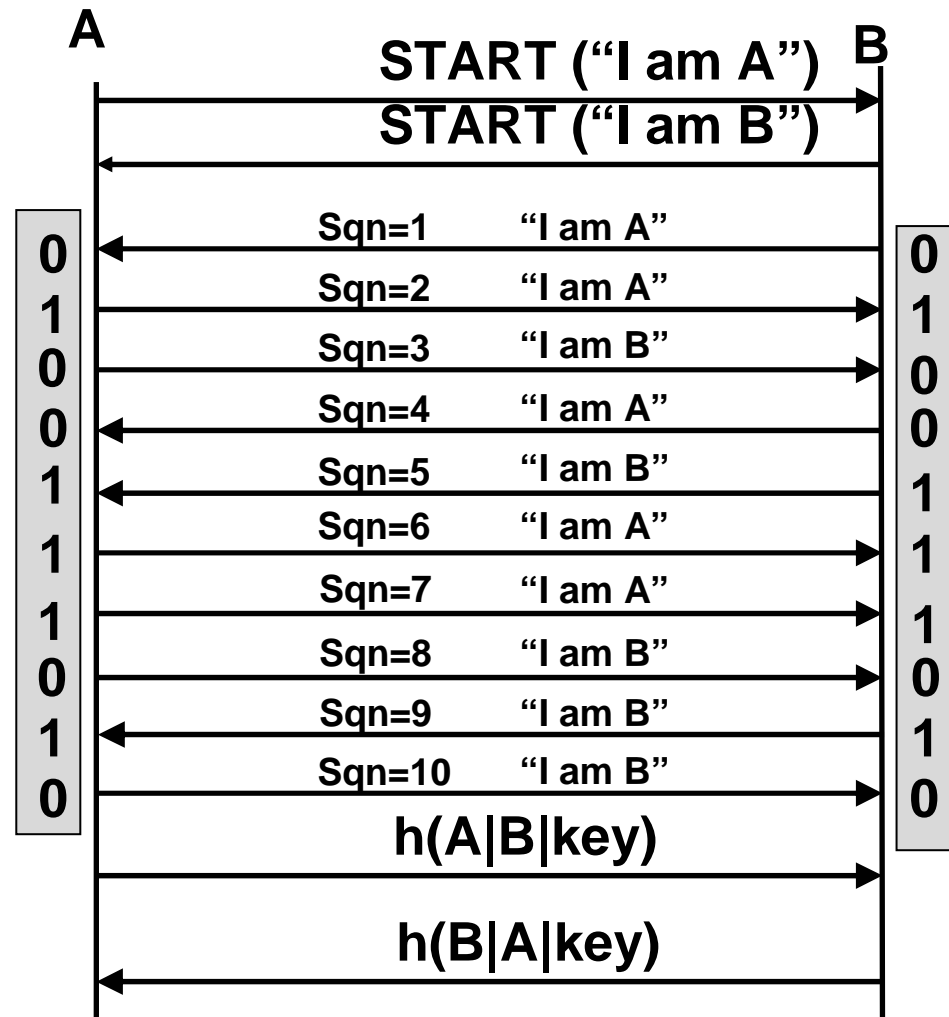
- If you move the devices at a high speed, the doppler effect might solve the problem for you ☺ !
- A more practical solution is to add a random frequency offset so that A and B span over similar frequency ranges.
 - Btw This solution does not require to modify the standard!

Packets belong to A or B!
We want to use these packets!



The Shake' em Up protocol (STU)

- We combine the previous protocol with shaking.
- A user that wants to pair to devices A and B
 - Brings the devices together
 - Shakes them up!
 - Triggers the protocol (for example by pushing a button on the devices)...



Performance: Energy Consumption

- In STU, each device
 - processes N small messages, where N is # of bits of the secret (total number of bits sent: 2016)
 - ...but performs almost no computation.
- In a DH based scheme,
 - each node sends only one large message (>1024 bits)...
 - but performs a lot of computation...i.e. 4.12×10^8 single precision multiplications (if $N=72$).
- By using the heuristic that transmitting one bit consumes as much energy as executing 800 instructions...
 - we can show (see paper) that our scheme is 100 times more energy efficient than a DH based scheme!
 - Elliptic Curve DH can reduce the communication cost by 5....but our scheme would still be more efficient.

Conclusion

- We've presented a key exchange protocol for CPU constrained devices that uses CSMA-based wireless communication.
- We believe this is the first solution that does not rely on cryptography, out-of-band channel or specialized hardware.
 - Very simple and efficient
 - Does not require computation...just transmission and good shaking!
- This is only one example of the ACI SPLASH results

Selected Results in 2005

- Key Establishment
 - ***Shake Them Up!***. ACM/Usenix *Mobisys'05*
 - ***Self-Configurable Key Pre-distribution Mobile Ad Hoc Networks, IFIP Networking'05.***
 - ***Authentication using ID-based Hash chains***, in submission
- Secure Aggregation
 - ***Efficient Aggregation of Encrypted Data in Wireless Sensor Networks***, ACM/IEEE *Mobiquitous* Conference
 - ***Secure Acknowledgment Aggregation, Computer Networks (Elsevier).***
- Cooperation Enforcement
 - ***Pocket Bluff***, in submission.
 - ***Real life experience of Cooperation Enforcement Based on Reputation (CORE) for MANETs, IEEE REALMAN***
 - ***Non cooperative forwarding in ad hoc networks, IFIP Networking'05***
 - ***Some game-theoretic problems in wireless ad hoc networks, NGI 2005***
 - ***Analysis of coalition formation and cooperation strategies in mobile ad hoc networks, Ad Hoc Networks Journal***
- Secure Routing
 - ***Securing Route Discovery in DSR, ACM/IEEE Mobiquitous'05***
 - ***Ad hoc networks security***, Chapter in the book "Handbook of information security"

Diffusion

- We organized:
 - **ESAS 2005** (2nd European Workshop on Security and Privacy in Ad hoc and Sensor Network)
 - R. Molva (Eurecom) and G. Tsudik (UCI) were co-chairs
 - C. Castelluccia was in the steering and PC committees.
 - First International Workshop on Trust, Security and Privacy for Ubiquitous Computing (**TSPUC 2005**)
 - R. Molva was co-chair
 - C.Castelluccia, P.Michiardi and G.Tsudik were in the PC
 - First IEEE Conference on Security and Privacy for Emerging Areas in Communication and Networks (**SecureComm'05**)
 - G.Tsudik was co-chair
 - C.Catelluccia, P. Michiardi and R.Molva were in the PC

Future Work

- INRIA will participate to the IST STREP project UbiSen&Sec (Starting 2006)
 - This project is about security in wireless sensor networks
- EURECOM is a partner of two IST FET Projects on Autonomic Computing: HAGGLE and CASCADAS
- We will organize:
 - ESAS'06 (together with ESORICS'06)
 - TSPUC'06
 - SecureComm'06
 - IEEE ICC 2006 Network Security and Information Assurance Symposium

Collaborations within Splash

- The collaboration within the project has been very active
 - C.Castelluccia (INRIA) has spent 2 years at UCI
 - P.Mutaf (INRIA) has moved to Sophia-Antipolis to closely collaborate with Eurecom
 - P. Michiardi, R. Molva (EURECOM) worked with E. Altman (INRIA) on Game theory (2 papers)
 - N.Saxena (UCI) is currently visiting INRIA for 3 months.